

Lecture 15 CPSC 447/547 - Intro to QC

Grover's Search Algorithm

Outline

- Unstructured search problem
 - Reflection Operators.
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So far in lectures, we've been using quantum algorithms that follow the **same recipe**:

- ① Uniform **Superposition**
- ② **Query** an oracle in superposition
- ③ Directly **measure** or transform-then-measure.

- Can we use a quantum oracle differently?
- What more tasks can we do if we use oracle repeatedly?

To answer those questions, we need to understand the role of an oracle better.

Phase Oracle (revisited)

$$O_f^\pm |x\rangle = (-1)^{f(x)} |x\rangle$$

$$O_f^\pm \left(\sum_x \alpha_x |x\rangle \right) = \sum_x \alpha_x (-1)^{f(x)} |x\rangle$$

Observation: "Negate α_x to $-\alpha_x$ if $f(x) = 1$ ".

What if: $f(x) = 1$ very rare.

E.g. $f(x) = \begin{cases} 1 & \text{if } x = z, \text{ for some } z \in \{0, 1\} \\ 0 & \text{otherwise} \end{cases}$.

"There exists a single input, s.t. f evaluates to 1".

What is O_f^\pm in this case?

$$O_f^\pm |x\rangle = \begin{cases} -|x\rangle & \text{if } |x\rangle = |z\rangle \\ |x\rangle & \text{otherwise} \end{cases}$$

We have seen this type of transformation: "**Reflection**!"

$$R_z^\pm = I - 2\pi_z, \text{ where } \pi_z = |z\rangle\langle z|$$

\uparrow reflection \uparrow projector on to $|z\rangle$

$$\Rightarrow (I - 2\pi_z) |z\rangle = |z\rangle - 2|z\rangle\langle z|z\rangle = -|z\rangle$$

hyperplane.

$$\begin{aligned} \text{and } (I - 2\pi_z) |x\rangle &= |x\rangle - 2|z\rangle\langle z|x\rangle^0 \quad \text{for } |x\rangle \perp |z\rangle \\ &= |x\rangle \end{aligned}$$

$\Rightarrow O_z^\pm = R_z^\pm$ = "Reflection through (hyperplane of) $|z\rangle$ ".

Example "Reflection through $|0^n\rangle$ ".

$$R_0^\pm |x\rangle = \begin{cases} -|x\rangle & \text{if } |x\rangle = |0^n\rangle \\ |x\rangle & \text{otherwise.} \end{cases}$$

$$|\psi\rangle = \sum_x \alpha_x |x\rangle \xrightarrow{R_0^\pm} |\phi\rangle = -\alpha_0 |0^n\rangle + \sum_{x \neq 0^n} \alpha_x |x\rangle.$$

Unitary matrix for R_0^\pm :

$$R_0^\pm = I - 2\pi_0 = I - 2|0^n\rangle\langle 0^n|.$$

Can we implement R_0^\pm with quantum circuit?

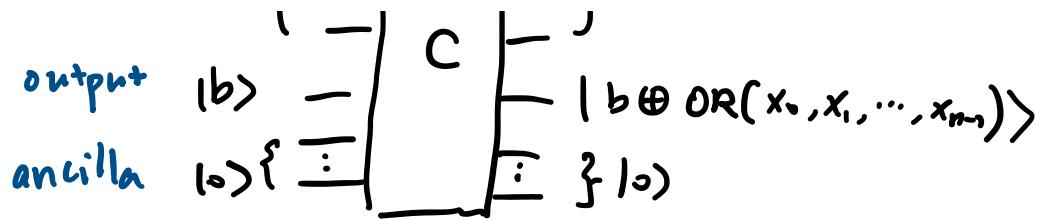
Circuit Implementation of Reflections

① Example R_0^\pm : Negate $|x\rangle$ if and only if $x = 0^n$
i.e., $\text{NOT}(\text{OR}(x_0, x_1, \dots, x_{n-1}))$

Using the Boolean circuit for OR function?

- OR function can be implemented by $O(n)$ Toffoli gates.
- Make it reversible.

Input $|x\rangle \xrightarrow{=} F |x\rangle$



- Build R_o^\pm : $|x\rangle \longrightarrow (-1)^{\text{OR}(x_0, \dots, x_{n-1})} |x\rangle$.

Let $|b\rangle = |-\rangle$, then

$$|x\rangle |b\rangle |0\rangle = |x\rangle |-\rangle |0\rangle \xrightarrow{C} (-1)^{\text{OR}(x_0, \dots, x_{n-1})} |x\rangle |-\rangle |0\rangle \quad (\text{smiley face})$$

$$= -R_o^\pm |x\rangle |-\rangle |0\rangle$$

So

$$|x\rangle \left\{ \begin{array}{c} \vdots \\ \boxed{R_o^\pm} \\ \vdots \end{array} \right\} \Rightarrow |x\rangle \left\{ \begin{array}{c} \vdots \\ \boxed{C} \\ \vdots \end{array} \right\} -R_o^\pm |x\rangle$$

$$\begin{array}{c} |-\rangle \\ \hline |0\rangle \end{array} \quad \begin{array}{c} |-\rangle \\ \hline |0\rangle \end{array}$$

② One more example : Reflection through $|+\rangle$.

$$R_+^\pm |x\rangle = \begin{cases} -|x\rangle & \text{if } |x\rangle = |+\rangle \\ |x\rangle & \text{otherwise.} \end{cases}$$

Unitary Matrix : $R_+^\pm = I - 2|+\rangle\langle +|$

Circuit Implementation (using R_o^\pm)

$$|x\rangle \left\{ \begin{array}{c} \vdots \\ \boxed{R_+^\pm} \\ \vdots \end{array} \right\} = \begin{array}{c} -|-\rangle \\ \vdots \\ -|H\rangle \end{array} \left\{ \begin{array}{c} \vdots \\ \boxed{R_o^\pm} \\ \vdots \end{array} \right\} \begin{array}{c} |H\rangle \\ \vdots \\ |H\rangle \end{array}$$

1. Change $|+\rangle$ to $|-\rangle$

- 1. Change it to $|0\rangle$,
- 2. Reflect through $|0^n\rangle$.
- 3. Change back.

Use Reflections for (unstructured) search (Grover)

Prob. Given an **oracle access** to function

$$f: \{0,1\}^n \rightarrow \{0,1\},$$

$$f(x) = \begin{cases} 1 & \text{if } x = z \text{ for some } z \in \{0,1\}^n \\ 0 & \text{otherwise} \end{cases}$$

Want to find z .

"Search for x s.t. $f(x) = 1$ ".

Quantum Strategy:

Design a quantum state $|+\rangle$ s.t.

when we measure, we obtain $|z\rangle$ with high prob.

Initial: $|+\rangle = \frac{1}{\sqrt{2^n}} \sum_x |x\rangle$ "**uniform superposition**"
 (like a uniform random guess)

Some quantum circuit

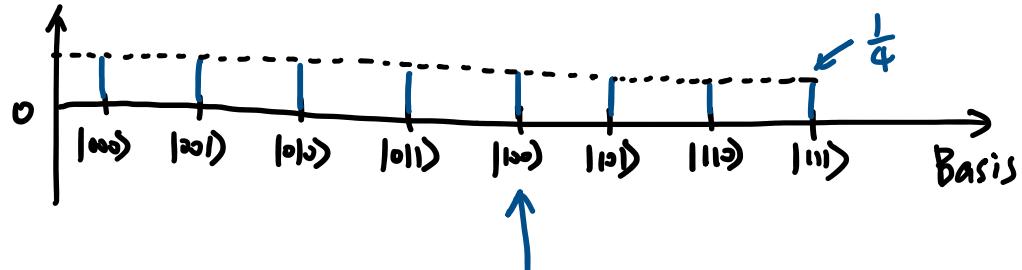
Final: $|+\rangle \approx |z\rangle$ "as **close to correct answer** z as possible".

Let's see how we can amplify the amplitude on $|z\rangle$.

Example : $n=3$. $f(x)=1$ if $x=|00\rangle$, and $f(x)=0$ otherwise,

$$\text{We start with } |\psi\rangle = \frac{1}{\sqrt{8}} \sum_{x \in \{0,1\}^3} |x\rangle = \frac{1}{\sqrt{8}} (|000\rangle + |001\rangle + |010\rangle + |011\rangle + |100\rangle + |101\rangle + |110\rangle + |111\rangle).$$

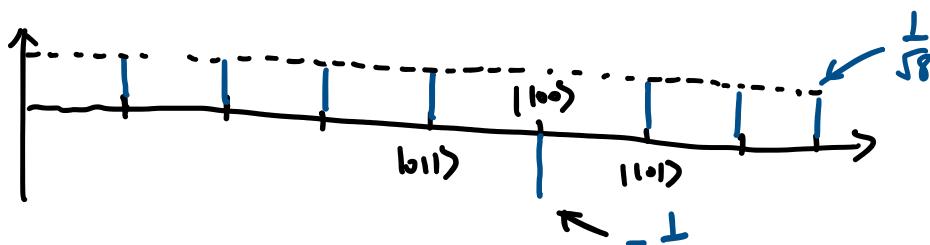
Amplitude of $|\psi\rangle$



Want to maximize amplitude here.

① Let's first try the phase oracle. $O_f^\pm = R_{|00\rangle}^\pm$

Amplitude of $O_f^\pm |\psi\rangle$



If we measure, $\text{Prob}(|100\rangle) = (-\frac{1}{\sqrt{8}})^2 = \frac{1}{8}$. ☺.

But at least we've "singled out" $|100\rangle$ after O_f^\pm .

② Keep going : Next try $R_+^\pm = I - 2|+\rangle\langle +|$

What's $(-R_+^\pm)O_f^\pm |\psi\rangle$? ↴ (Actually, need $-R_+^\pm$)

$$-R_+^\pm O_f^\pm |\psi\rangle = \underbrace{\left(2|+\rangle\langle +| - I \right)}_{\text{What's } (-R_+^\pm)O_f^\pm |\psi\rangle?} \underbrace{\left(\frac{1}{4} \sum_{x \neq |00\rangle} |x\rangle - \frac{1}{4} |100\rangle \right)}_{\text{Actually, need } -R_+^\pm}$$

$$\langle +^3 | \left(\sum_x \alpha_x |x\rangle \right) = \left(\frac{1}{\sqrt{2^n}} \sum_y \langle y | \right) \left(\sum_x \alpha_x |x\rangle \right)$$

$$= \frac{1}{\sqrt{2^n}} \sum_y \sum_x \alpha_x \underbrace{\langle y | x \rangle}_{\delta_{xy}}$$

$$= \frac{\sum_x \alpha_x}{\sqrt{2^n}} = \sqrt{2^n} \cdot M, \quad M = \frac{\sum_x \alpha_x}{2^n}.$$

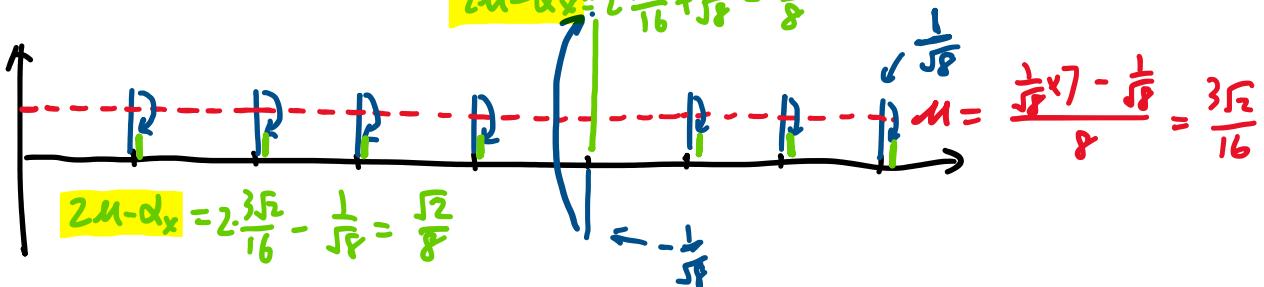
↑ Mean of amplitudes

$$\text{So, } (-R_+^\pm) \left(\sum_x \alpha_x |x\rangle \right)$$

$$= 2 \cdot \sqrt{2^n} M |+^3\rangle - \sum_x \alpha_x |x\rangle$$

$$= \sum_x (2M - \alpha_x) |x\rangle \text{ "Flip around mean":}$$

$$2M - \alpha_x = 2 \cdot \frac{3\sqrt{2}}{16} + \frac{1}{\sqrt{2}} = \frac{5\sqrt{2}}{8}$$



Before: $\frac{1}{\sqrt{8}}|000\rangle + \frac{1}{\sqrt{8}}|001\rangle + \frac{1}{\sqrt{8}}|010\rangle + \frac{1}{\sqrt{8}}|011\rangle - \frac{1}{\sqrt{8}}|100\rangle + \frac{1}{\sqrt{8}}|101\rangle + \dots + \frac{1}{\sqrt{8}}|111\rangle$.

After: $\frac{\sqrt{2}}{8}|000\rangle + \frac{\sqrt{2}}{8}|001\rangle + \frac{\sqrt{2}}{8}|010\rangle + \frac{\sqrt{2}}{8}|011\rangle + \frac{5\sqrt{2}}{8}|100\rangle + \frac{\sqrt{2}}{8}|101\rangle + \dots + \frac{\sqrt{2}}{8}|111\rangle$.

This is better! Amp on $|100\rangle$ is amplified. ☺

Next lecture: Can we amplify more?

What if we keep going: $(-R_+^\pm) O_f^\pm$?

How many iterations until the prob
of obtaining 100 is high enough?

“I am a man of my word;